

# **Compact Formulations of Network Flow Problems**

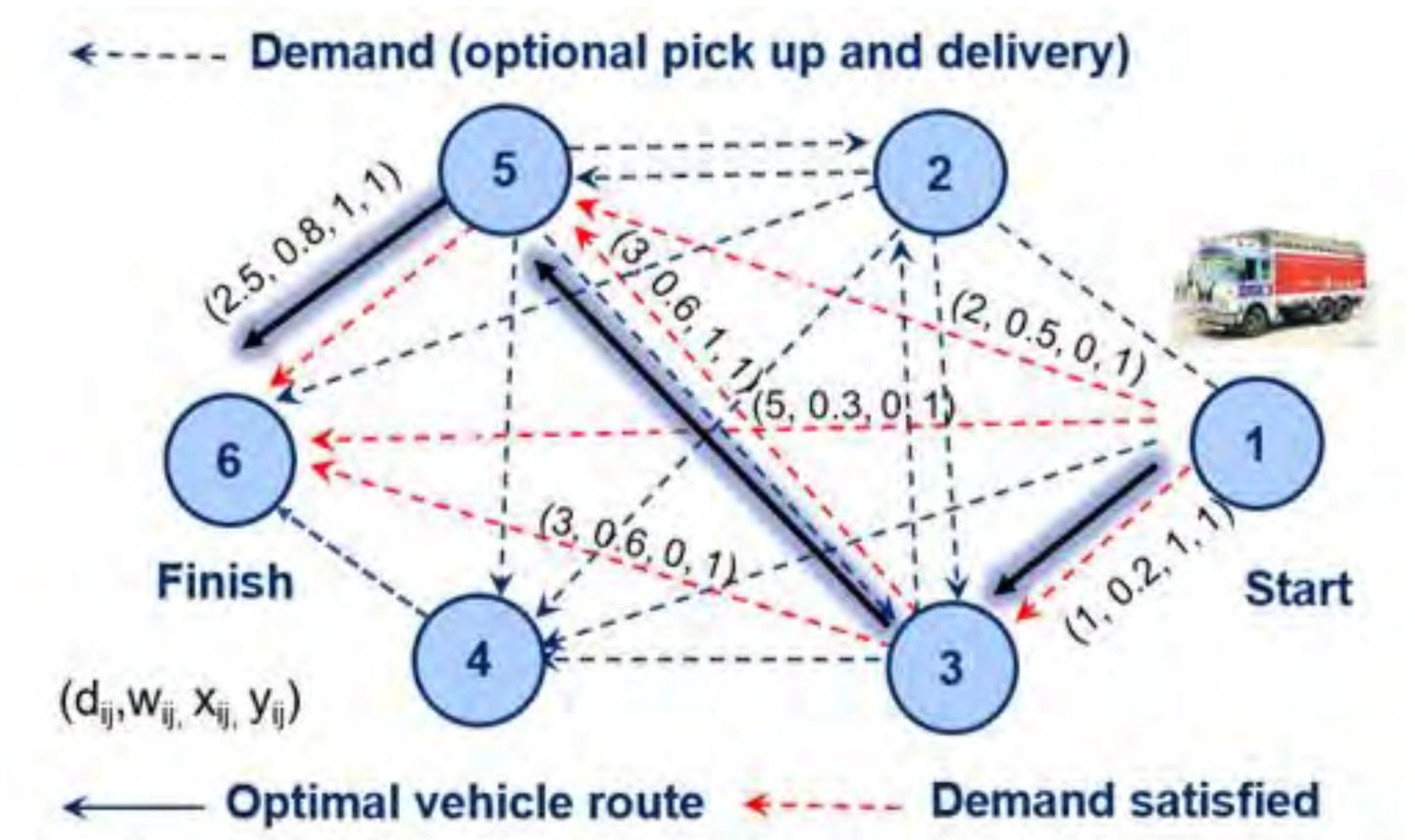
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## Abstract

The Triples formulation is a compact formulation of multicommodity network flow problems that provides a different representation of flow than the traditional widely used node-arc and arc-path approaches. We have applied the Triples concept to the backhaul profit maximization problem (BPMP) and the fixed charge network flow problem (FCNF). Experimental analysis shows that the Triples formulations of BPMP and FCNF are more compact, stronger, and faster than the traditional node-arc models for most problem instances.

### Backhaul Profit Maximization Problem



dig distance from i to j; will demand from i to j; x<sub>ii</sub>:=1 if the vehicle travels on arc (i,j); y ij: =1 if demand from i to j is satisfied

Objective: maximize profit Decision variables: x<sub>ii</sub> , y<sub>ii</sub> Constraints: time and weight.

# Flow: Node-Arc vs. Triples

#### Node-Arc Model

- $-z_{kl.ii}$  = 1 if demand from k to l is shipped on arc (i, j)
- -Example solution (6 demands selected):

•
$$z_{13,13} = 1$$
,  $z_{15,13} = z_{15,35} = 1$ ,  $z_{16,13} = z_{16,35} = z_{16,56} = 1$ 

$$z_{35,35} = 1$$
,  $z_{36,35} = z_{36,56} = 1$ ,  $z_{56,56} = 1$ 

 $-\theta_{ii}$  = total flow (tons) on arc (i, j)

$$\theta_{13} = w_{13} z_{13,13} + w_{15} z_{15,13} + w_{16} z_{16,13} = 1 \text{ ton}$$

•
$$\theta_{35}$$
= $w_{15} z_{15,35}$ +  $w_{16} z_{16,35}$ + $w_{35} z_{35,35}$ +  $w_{36} z_{36,35}$ = 2 tons

$$\bullet \theta_{56} = w_{16} z_{1656} + w_{36} z_{3656} + w_{56} z_{5656} = 1.7 \text{ ton}$$

- • $\theta_{56} = w_{16} z_{16,56} + w_{36} z_{36,56} + w_{56} z_{56,56} = 1.7 \text{ tons}$
- -O(n<sup>4</sup>) binary variables and O(n<sup>3</sup>) constraints

# Triples Model

- $-u_{ii}^{k}$  = flow (tons) from i to j diverted through k: routed on arc (i, k) and a path from k to j
- -Example solution diverts 3 of 6 selected demands

$$\mathbf{u}_{15}^3 = \mathbf{w}_{15}$$
,  $\mathbf{u}_{16}^3 = \mathbf{w}_{16}$ , and  $\mathbf{u}_{36}^5 = \mathbf{w}_{16} + \mathbf{w}_{36}$ 

$$\theta_{13} = w_{13} + u_{15}^3 + u_{16}^3 = 1 \text{ ton}$$

$$\theta_{35} = w_{35} + u_{15}^3 + u_{36}^5 = 2 \text{ tons}$$

 $\theta_{56} = w_{56} + u_{36}^{5} = 1.7 \text{ tons}$ 

-O(n<sup>3</sup>) binary variables and O(n<sup>2</sup>) constraints

# BPMP: Node-Arc vs. Triples Models

Maximize  $p\left[\sum_{(k,l)\in E}d_{kl}w_{kl}y_{kl}\right]-c\sum_{(i,j)\in E}\theta_{ij}d_{ij}-cv\sum_{(i,j)\in E}d_{ij}x_{ij}$ Subject to:

$$\sum_{j\in V} z_{kl,kj} = y_{kl} (k,l) \in A$$

$$\sum_{i \in V} z_{kl,il} = y_{kl} (k,l) \in A$$

$$\sum_{i \in V, (i,a) \in A} z_{kl,ia} = \sum_{j \in V, (a,j) \in A} z_{kl,aj} \quad (k,l) \in A, a \in V \setminus \{k,l\}$$

$$\sum_{(k,l)\in A} z_{kl,ij} \leq (n^2 - 3n + 3)x_{ij} \ (i,j) \in A$$

$$\sum_{(\mathbf{1},j)\in A} x_{\mathbf{1},j} = \mathbf{1}$$

$$\sum_{(i,n)\in A} x_{i,n} = 1$$

$$\sum_{i\in V\setminus\{k,n\}}x_{ik}=\sum_{j\in V\setminus\{1,k\}}x_{kj} \quad k\in V\setminus\{1,n\}$$

$$\sum_{(i,j)\in A} d_{ij} x_{ij} \leq D$$

$$\theta_{ij} = \sum_{(\mathbf{k},\mathbf{l})\in A} \mathbf{w}_{\mathbf{kl}} \mathbf{z}_{\mathbf{kl},ij} \quad (\mathbf{i},\mathbf{j})\in A$$

$$\theta_{ij} = w_{ij}y_{ij} + \sum_{(i,k,j)\in T} u_{ik}^{j} + \sum_{(k,j,i)\in T} u_{kj}^{i} - \sum_{(i,j,k)\in T} u_{ij}^{k}, (i,j) \in A$$

$$\theta_{ij} \leq Q, (i,j) \in A$$

$$\theta_{ij} \leq Qx_{ij}, (i,j) \in A$$

$$s_i - s_j + (n)x_{ij} \le n-1 \ (i,j) \in A$$

$$x_{ij} = 0, 1 \ (i,j) \in A$$

$$y_{kl} = 0, 1 \ (k, l) \in A$$

$$u_{ij}^k \geq 0$$
,  $(i,j,k) \in T$ 

$$z_{kl,ij} = 0, 1 \ (k,l) \in A, (i,j) \in A$$

Black: common for both models

Blue: unique to Node-arc model

Red: unique to Triples model

# Restricted Triples Heuristic

Algorithm: Restricted Triples Heuristic Input: A BPMP instance

Output: A feasible triples solution

#### (1) $\hat{\mathcal{T}} \leftarrow \emptyset$

(2) For 
$$\forall (i,j,k) \in \mathcal{T}$$

$$\rho_{ij}^k \leftarrow p d_{ij} w_{ij} - c(d_{ik} + d_{kj})(v + w_{ij})$$

4) If 
$$w_{ij} + w_{ik} \le Q$$
 Then  $\rho_{ij}^k \leftarrow \rho_{ij}^k + (p-c)d_{ik}w_{ik}$ 

(5) If 
$$w_{ij} + w_{kj} \le Q$$
 Then  $\rho_{ij}^k \leftarrow \rho_{ij}^k + (p-c)d_{kj}w_{kj}$ 

(i) If 
$$\rho_{ij}^k \ge 0$$
 Then  $\hat{\mathcal{T}} \leftarrow \hat{\mathcal{T}} \cup \{(i,j,k)\}$ 

(7) End For

(8) Solve enhanced triples formulation with  $\hat{\mathcal{T}}$  (i.e., fix  $u_{ij}^k = 0$  for  $(i, j, k) \in \mathcal{T} \setminus \hat{\mathcal{T}}$ )  $(9) \mathcal{A}_x \leftarrow \{(i,j) \in \mathcal{A} : x_{ij} = 1\}$ 

#### (10) $\mathcal{R}_y \leftarrow \{(i,j) \in \mathcal{R} : y_{ij} = 1\}$

(11) Solve enhanced triples formulation with  $\mathcal{T}$  subject to  $x_{ij} = 1 \quad \forall (i,j) \in A_x,$ 

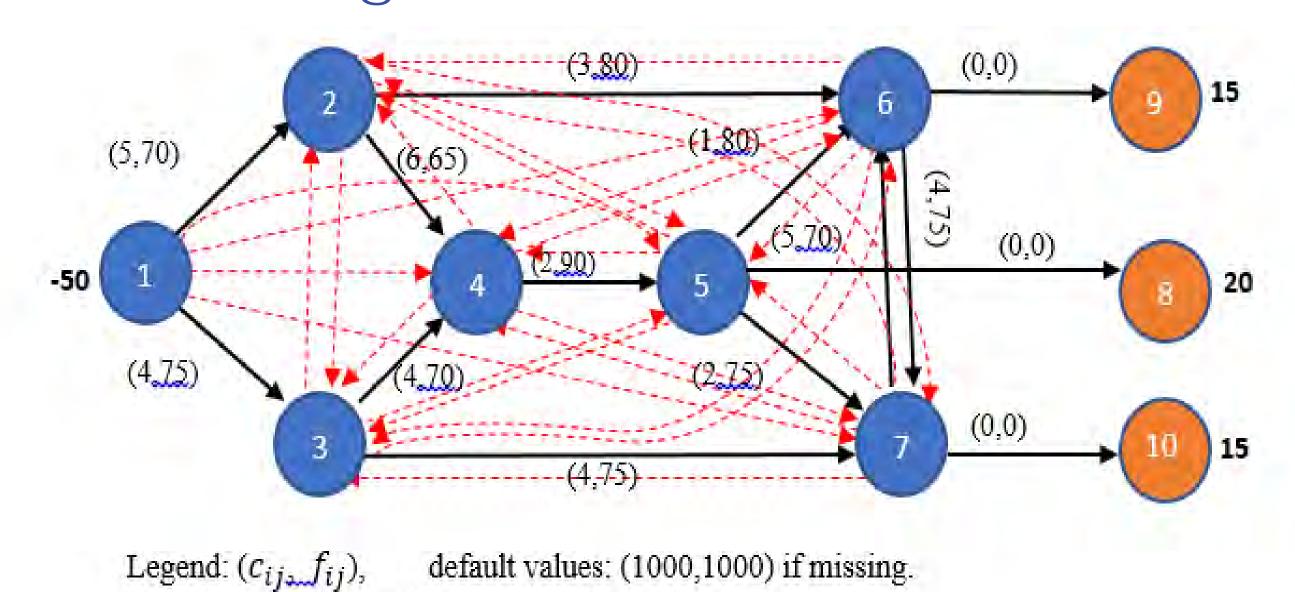
 $y_{ij} = 1 \quad \forall (i,j) \in R_y.$ 

(12) **Return** Triples solution (x, y, u, s)

# Empirical CPU Comparison

	Average CPU (seconds)						
Nodes	Node Arc	Triples	Speed Up				
10	14.4	2.4	1.5				
20	13,644	20	<b>758</b>				
30	N/A	480	NA				

# Fixed Charge Network Flow Problem



#### Base

$$\begin{aligned} & \text{Min } \sum_{(i,j) \in A} c_{ij} \, x_{ij} + \sum_{(i,j) \in A} f_{ij} y_{ij} \\ & \sum_{(i,k) \in A} x_{ik} - \sum_{(k,j) \in A} x_{kj} = b_k, \ k \in V \\ & x_{ij} \leq (\sum_{\mathbf{d} \in D} b_{\mathbf{d}}) y_{ij}, \, (i,j) \in A \\ & x_{ij} \geq 0, \, (i,j) \in A \\ & y_{ij} \in \{0,1\} \end{aligned}$$

#### MCE

$$\begin{aligned} & \text{Min } \sum_{(i,j) \in A} x_{ij} c_{ij} + \sum_{(i,j) \in A} f_{ij} y_{ij} \\ & \sum_{(i,k) \in A} w_{ik}^d - \sum_{(k,j) \in A} w_{kj}^d = \begin{cases} b_d & \text{k} = \text{d} \in D \\ 0, & \text{k} \in \mathbb{V} \backslash 1, \text{d} \in D, d \neq k \\ -b_d & \text{k} = 1, \text{d} \in D \end{cases} \\ & x_{ij} \leq (\sum_{\text{d} \in D} b_d) y_{ij}, (i,j) \in A \\ & x_{ij} \geq \sum_{\text{d} \in D} w_{ij}^d (i,j) \in A \\ & w_{ij}^d \leq b_d y_{ij}, (i,j) \in A, d \in D \\ & w_{ij}^d \geq 0, (i,j) \in A, d \in D \\ & y_{ij} \in \{0,1\} \end{aligned}$$

# **Triples**

Minimize  $\sum_{(i,j)\in A} x_{ij} c_{ij} + \sum_{(i,j)\in A} f_{ij} y_{ij}$  $x_{ij} = d_{ij} + \sum_{(i,k,j) \in T} u_{ik}^j + \sum_{(k,j,i) \in T} u_{kj}^i - \sum_{(i,j,k) \in T} u_{ij}^k \ (i, j) : i \in V, j \in V \setminus \{i\}.$  $T: \{(i, j, k): i \in V, j \in V \setminus \{i\}, k \in V \setminus \{i, j\}\}$  $x_{ij} \leq s_{ij} * y_{ij}, (i, j) : i \in \mathbb{V}, j \in \mathbb{V} \setminus \{i\}$  $u_{ij}^k \ge 0 \quad \forall (i,j,k) \in T$  $y_{ij} \in \{0,1\} \quad \forall (i,j) \in A$ 

# Experimental Comparison

	MCE vs. Base			Triples vs. Base				
Node	CPU	ticks	Real	LP lower	CPU	ticks	Real	LP lower
Node	speed	speed	time	improve	speed	speedup	time	improve
	up	up	speedup	ment	up	эресиир	speedup	ment
20	6.45	1.34	1.91	36.8%	6.41	1.34	1.94	36.8%
30	6.15	1.12	1.56	26.6%	5.85	1.12	1.51	26.6%
100	7.02	1.44	1.62	89.1%	6.79	1.44	1.59	89.1%
500	1.41	16.29	0.12	320.1%	1.38	16.29	0.11	320.1%
Average	5.26	5.05	1.30	118.1%	5.10	5.05	1.29	118.1%

# Conclusions and Future Work

- The Triples model is more compact and faster to solve than node-arc model for BPMP.
- The Triples model is at least as good as the traditional advanced MCE formulation for FCNF.
- In the future, we will explore singleton-selection strategies to further accelerate our solution approach.