Memory-efficient approximation algorithms for Max-Cut and Max-k-Cut

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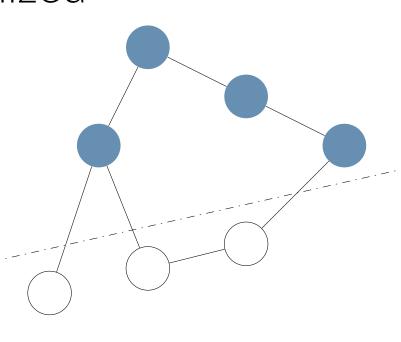
²Indian Institute of Technology Bombay

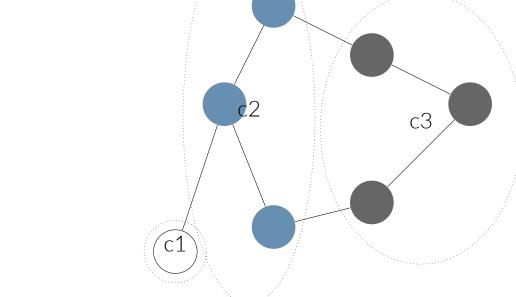
³Monash University

MAXCUT and MAX-K-CUT

MaxCut Partition nodes into two sets so that the weights of the edges crossing the partition is maximized

Max-k-Cut Partition nodes into k sets so that the weights of the edges crossing the partition is maximized





SDP Relaxation:

$$\max \langle cL_G, X \rangle$$

s.t. diag $(X) = 1$, $X \succeq 0$

SDP relaxation:

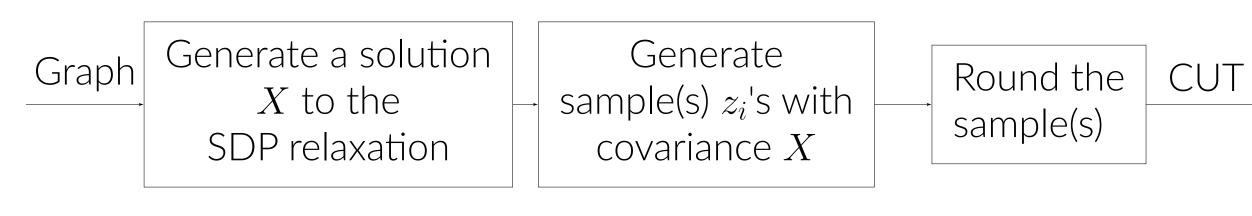
$$\max \quad \langle c_k L_G, X \rangle$$
s.t.
$$X_{ij} \geq -\frac{1}{k-1} \quad i \neq j,$$

$$\operatorname{diag}(X) = 1, X \succeq 0$$

Decision variable X uses n^2 storage! Too large for large-scale problems

First aim: Provide a low-memory representation of the decision variable

SDP-based rounding method



Generation of a cut only requires $z \sim \mathcal{N}(0, X)!$

If opt is the optimal cut value

Goemans-Williamson (GW) rounding: Frieze-Jerrum (FJ) rounding:

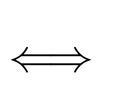
$$\mathbb{E}[\mathsf{CUT}] \geq \alpha_{GW}\mathsf{opt}[2]$$

$$\mathbb{E}[\mathsf{k}\text{-}\mathsf{CUT}] \geq \alpha_k \mathsf{opt}[1]$$

Second aim: Provide a low-memory implementation of the GW and FJ rounding that preserves approximation guarantees

Gaussian Sampling: A way to achieve our first aim

Nonnegative weighted sum of samples



Nonnegative weighted sum of matrices

- Let $z_t \sim \mathcal{N}(0, X_t)$ and $h_t \sim \mathcal{N}(0, H_t)$
- If $z_{t+1} = \sqrt{1-\gamma_t} z_t + \sqrt{\gamma_t \alpha} \zeta h_t$ for $\zeta \sim \mathcal{N}(0,1)$, then

$$\mathbb{E}\left[z_{t+1}z_{t+1}^{T}\right] = (1 - \gamma_t)X_t + \gamma_t H_t = X_{t+1}$$

Frank-Wolfe with Gausssian Sampling (FW-GS) [3]

Frank-Wolfe update step

$$X_{t+1} = (1 - \gamma)X_t + \gamma H_t$$

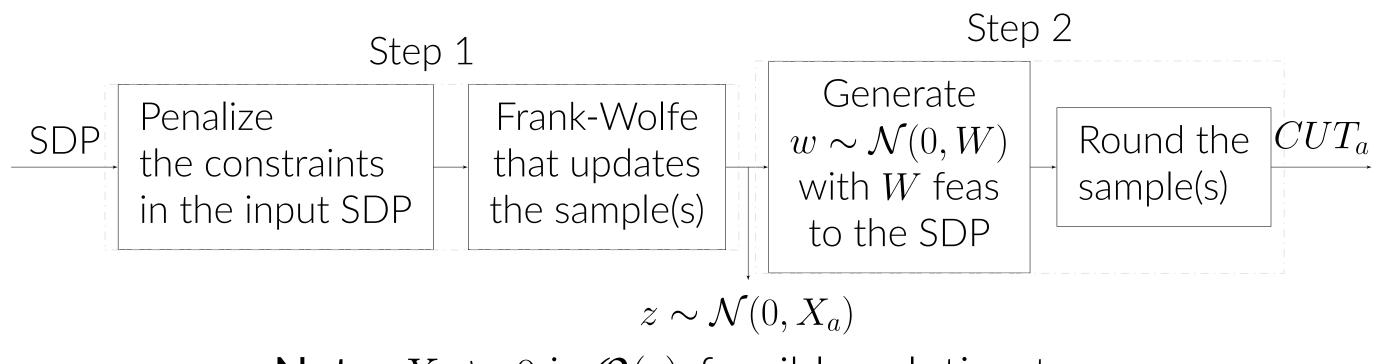
• Rank of update direction H_t at most one for the problem

$$\max_{X: \operatorname{Tr}(X) \leq \alpha, X \succeq 0} \quad g([\langle B_i, X \rangle]_{i=1}^d)$$

Not necessary to compute X_{t+1} explicitly if we use Gaussian sampling!

Gaussian sample z: a low-memory representation of X

Low-memory implementation of GW rounding: Our second aim (Part 1)



Note: $X_a \succeq 0$ is $\mathcal{O}(\epsilon)$ -feasible solution to input SDP with d linear constraints

Key features: Uses $\mathcal{O}(n+d)$ memory, preserves convergence rate of Frank-Wolfe

MaxCut using FW-GS: $\mathbb{E}[\mathsf{CUT}_a] \geq \alpha_{GW}(1-2\epsilon)\mathsf{opt}$

What about Max-k-Cut? SDP relaxation of Max-k-Cut has n^2 constraints. FW-GS uses $\mathcal{O}(n^2)$ memory!

Low-memory implementation of FJ rounding: Our second aim (Part 2)

A new relaxation of Max-k-Cut

 $\max \quad \langle c_k L_G, X \rangle$

s.t. $X_{ij} \geq -\frac{1}{k-1}$ $(i,j) \in E, i < j,$ diag $(X) = 1, X \succeq 0$

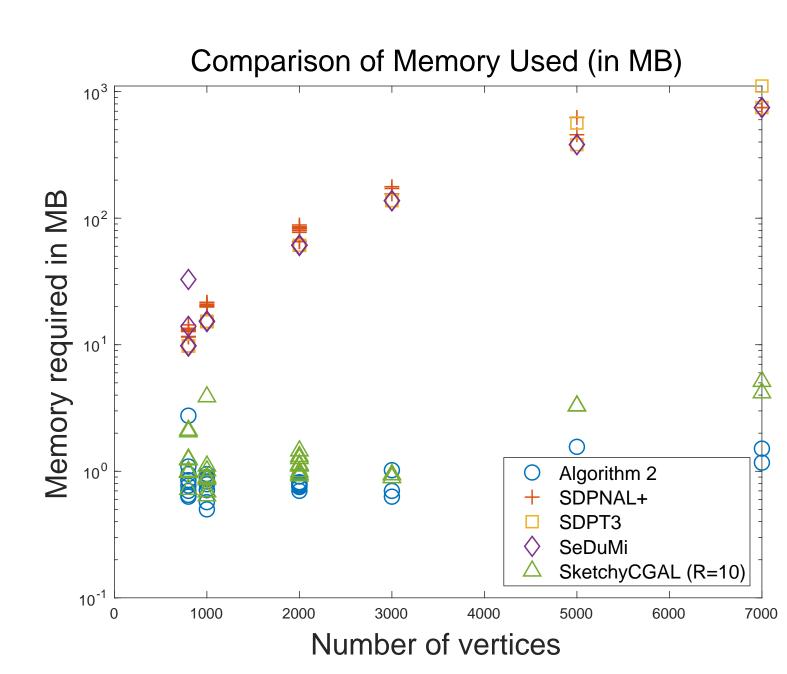
Key observation: If the rounding is applied to the new relaxation

 $\mathbb{E}[\mathsf{k}\text{-}\mathsf{CUT}] \geq \alpha_k \mathsf{opt}$

Max-k-Cut using FW-GS: $\mathbb{E}[\mathsf{CUT}_a] \geq \alpha_k (1 - 4\epsilon) \mathsf{opt}$

Implementing FW-GS for MaxCut with $\epsilon=0.1$

Specifications: Used a machine with 8GB RAM and 4 cores



- Simple to implement
- Offers scope for improvement in computational time
- Memory used is linear in n and significantly lower than standard SDP solvers

Future Directions

- What if the input graph for Max-k-Cut is dense? Sparsification
 of the graph to reduce the number of edges
- Are there other problems which can be solved using FW-GS?

References

- [1] Alan Frieze and Mark Jerrum. Improved approximation algorithms for maxk-cut and max bisection. *Algorithmica*, 18(1):67--81, 1997.
- [2] Michel X Goemans and David P Williamson. Improved approximation algorithms for maximum cut and satisfiability problems using semidefinite programming. *Journal of the ACM (JACM)*, 42(6):1115--1145, 1995.
- [3] Nimita Shinde, Vishnu Narayanan, and James Saunderson. Memory-efficient structured convex optimization via extreme point sampling. arXiv preprint arXiv:2006.10945, 2020.